

Enumerability & diagonalization

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Exercise

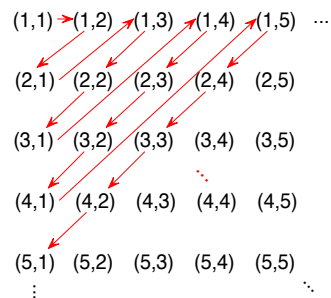
Show the following statements.

1. Every finite set is enumerable.
2. If a **non-empty** set A is enumerable, then it is enumerable by a **total** function. Why does this not work for the empty set?

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Cantor's Zig-Zag

Pairs of integers: Cantor's Zig-Zag.



$$f(1) = (1, 1), f(2) = (1, 2), f(3) = (2, 1), f(4) = (1, 3), f(5) = (2, 2), \dots$$

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Code numbers

Definition. Given an enumeration f of a set A , a **code numbers** of an element a of A is a number n such that $f(n) = a$.

Examples:

- The code number of $(1, 3)$ with respect to Cantor's Zig-Zag is 4.
- For the enumeration $2, 2, 4, 4, 6, 6, \dots$ of the even numbers, the code numbers of 4 are 2 and 3.

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Encodings

Definition. An **encoding** of a set A is a total injective function $c : A \rightarrow N$ into the natural numbers. For a in A , the number $c(a)$ is called the **code** of a .

Proposition. A set A has an encoding if and only if it is enumerable.

The proof of this proposition follows on the next two slides.

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From enumeration to encoding

Let $f : N \rightarrow A$ be an enumeration. An encoding $c : A \rightarrow N$ is given by

$$c(a) = \text{some } n \text{ such that } f(n) \text{ is equal to } a.$$

Because we need the function c to be an encoding, it must be total and injective. It is total because f is surjective. It is injective because f cannot send some n to two different values a and a' .

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From encoding to enumeration

Let $c : A \rightarrow N$ be an encoding. Then an enumeration $f : N \rightarrow A$ is given by

$$f(n) = \begin{cases} a & \text{if } n = c(a) \\ \text{undefined} & \text{otherwise} \end{cases}$$

f is well-defined (i.e. there is only one a for every n) because c is injective.

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“Enumerable” vs. “equinumerous”

Proposition. Every enumerable set A is either finite or equinumerous with N .

Proof. Let A be a set which is enumerable but not finite. Let $c : A \rightarrow N$ be an encoding of A . We define a bijection $b : N \rightarrow A$ by

$b(1) =$ the element of A with the smallest code

$b(2) =$ the element of A with the 2nd smallest code

$b(3) =$ the element of A with the 3rd smallest code

\vdots

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“Enumerable” vs. “equinumerous”

Proposition. A set A is enumerable if and only if it is finite or equinumerous with N .

Proof. The left-to-right direction is the result on the previous slide. For the right-to-left direction, suppose first that A is equinumerous with N . Then there is a bijection $b : N \rightarrow A$. In particular, A is the range of b , so A is enumerable. Second, suppose that A is finite. Then by an earlier exercise A is enumerable.

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Encoding other kinds of pairs

Proposition. If A and B are enumerable sets, then so is the set $A \times B$ of pairs.

Proof. Let $c_A : A \rightarrow N$ and $c_B : B \rightarrow N$ be encodings of A and B . An encoding $c_{A \times B} : A \times B \rightarrow N$ is given by

$$c_{A \times B}(a, b) = c_{N \times N}(c_A(a), c_B(b)),$$

where $c_{N \times N}$ is some encoding of pairs of integers. ($c_{A \times B}$ is injective because c_A , c_B , and $c_{N \times N}$ are.)

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Encoding pairs of integers

As we have seen, Cantor’s Zig-Zag provides an encoding of pairs of natural numbers.

Here is an alternative encoding:

$$c_{N \times N}(m, n) = p^m \cdot q^n$$

where p and q are different primes. The total function $c_{N \times N}$ is injective owing to the **uniqueness of prime decomposition**.

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Exercises

Show the following statements:

1. If A , B , and C are enumerable sets, then the set of triples

$$A \times B \times C = \{(a, b, c) : a \in A, b \in B, c \in C\}$$

is enumerable.

2. If A_1, A_2, \dots, A_k are enumerable sets, then the set of k -tuples $A_1 \times A_2 \times \dots \times A_k$ is enumerable.

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Exercises

Show the following statements:

1. If A is enumerable and there is a surjective function $A \rightarrow B$, then B is enumerable.
2. If B is enumerable and there is a total injective function $A \rightarrow B$, then A is enumerable.

Remark: these two statements are useful for some of the following exercises.

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Exercises

Show that the following statements:

1. The set $P_{fin}(N)$ of **finite subsets** of N is enumerable.
2. The set $P_{fin}(A)$ of finite subsets of an arbitrary enumerable set A is enumerable.

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Exercises

Show that the following sets are enumerable:

1. The set Q^+ of positive rational numbers.
2. The set Q of all rational numbers.
3. The set $A \cup B$ for enumerable sets A and B .
4. The set A^* of strings over an enumerable alphabet A .

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